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# Closed timelike curves in measurement-based quantum computation

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[arXiv:1003.4971v1 \[quant-ph\]](https://arxiv.org/abs/1003.4971v1)

# Outline

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- Closed time-like curves (CTCs) = time travel
- Deutsch's model for CTCs
- How CTCs show up in measurement-based quantum computation
- CTC model by Bennett/Schumacher/Svetlichny
- Conclusion

# Time travel

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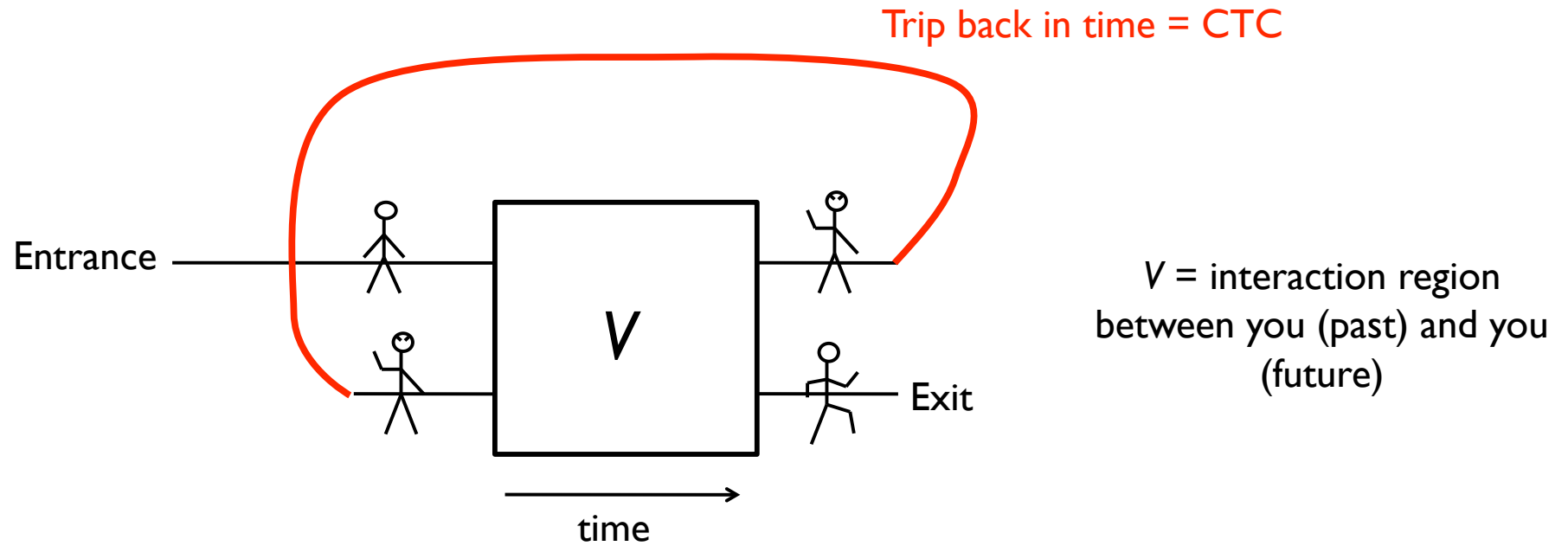
- To the future? Easy, use relativity.



- To the past? More involved:
  - Relativity predicts solutions with *closed timelike curves* (CTCs)= travel to the past
  - It's not known whether CTCs are physically possible, perhaps with clever black-hole engineering.
  - To avoid paradoxes, self-consistency conditions on time-travellers must apply – more on that later.

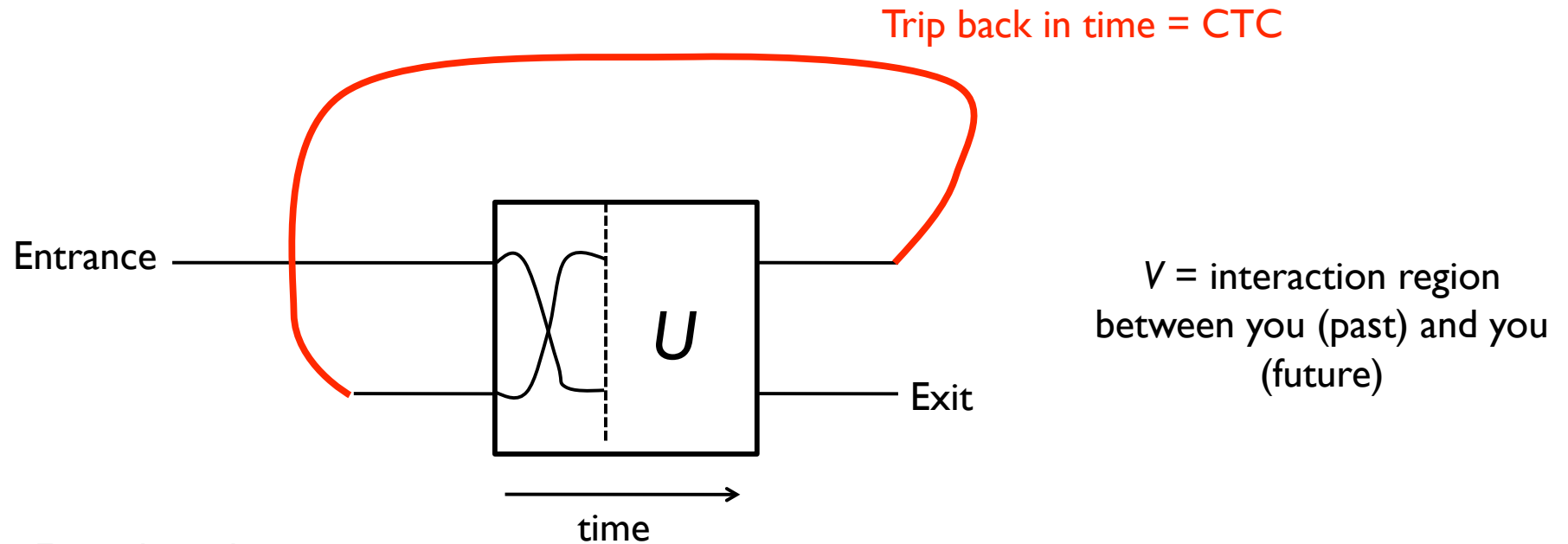
# Time travel - CTCs

- Time travel scenario:

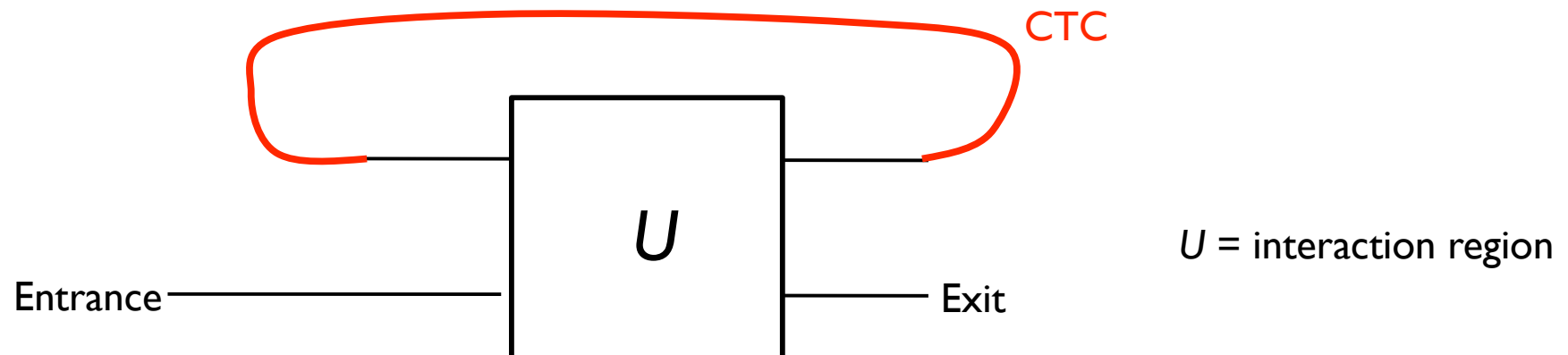


# Time travel - CTCs

- Time travel scenario:



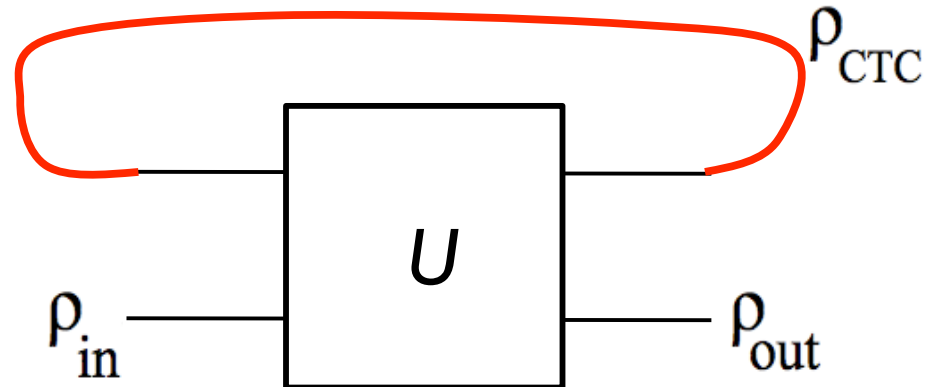
- Equivalent alternative:



# Deutsch's model for CTCs

Deutsch, Phys. Rev. D 44, 3197 (1991)

- $U = 2$ -qubit unitary
  - 1 qubit travels back in time
  - 1 qubit doesn't



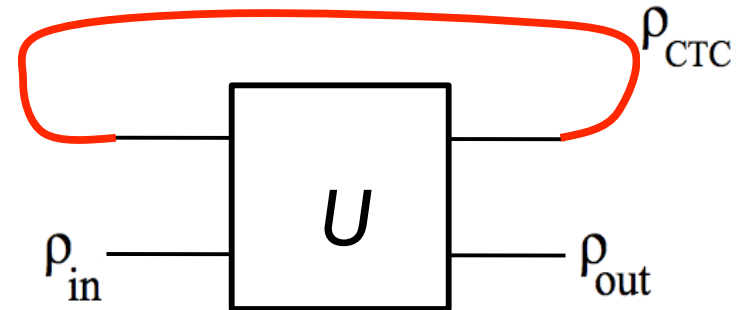
- Initial state:  $\rho_{CTC} \otimes \rho_{in}$
- After  $U$ :  $U(\rho_{CTC} \otimes \rho_{in})U^\dagger$
- Self-consistency condition:  $\rho_{CTC} = \text{Tr}_B \left[ U(\rho_{CTC} \otimes \rho_{in})U^\dagger \right]$
- Deutsch showed that:
  - there's always at least 1 self-consistent solution;
  - there can be multiple such solutions;
  - each solution corresponds to an input-output map, in general non-linear.

# Deutsch's model for CTCs

Deutsch, Phys. Rev. D 44, 3197 (1991)

- Some characteristics:

- avoids paradoxes;
- Under-determination of multiple solutions/maps: maxent? wormhole initial conditions? ...

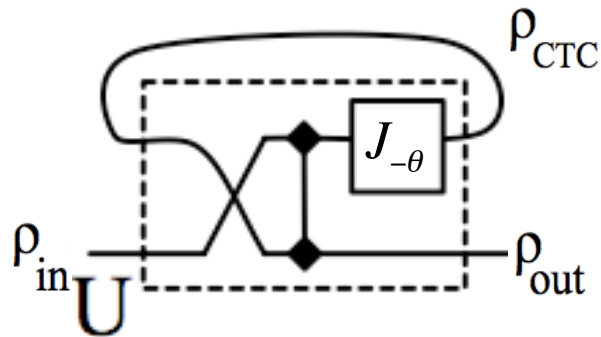


- **Computational power** of Deutsch's CTCs :

Non-orthogonal state discrimination? Solution to NP-complete problems?

- Bacon [arXiv:quant-ph/0309189v3](https://arxiv.org/abs/quant-ph/0309189v3) (solves NP-complete problems)
- Brun *et al.* [arXiv:0811.1209v2](https://arxiv.org/abs/0811.1209v2) (non-orthogonal state discrimination)
- Aaronson, Watrous: [arXiv:0808.2669v1](https://arxiv.org/abs/0808.2669v1) (CTCs  $\rightarrow$  PSPACE)
- Bennett *et al.* [arXiv:0908.3023v2](https://arxiv.org/abs/0908.3023v2) (criticism to results above)
- Cavalcanti, Menicucci [arXiv:1004.1219v2](https://arxiv.org/abs/1004.1219v2) (criticism of the criticism...)

# Deutsch - example



$$= CTR - Z = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix}$$

$$J_{-\theta} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & e^{-i\theta} \\ 1 & -e^{-i\theta} \end{pmatrix}$$

$$\rho_{in} = \frac{1}{2}(1 + \vec{n} \cdot \vec{\sigma}) \quad \rho_{CTC} = \frac{1}{2}(1 + \vec{m} \cdot \vec{\sigma}) \quad \rho_{out} = \frac{1}{2}(1 + \vec{r} \cdot \vec{\sigma})$$

- Self-consistency:

$$\begin{aligned} m_x &= n_z, \\ m_y &= m_z(n_x \sin \theta - n_y \cos \theta), \\ m_z &= m_z(n_x \cos \theta + n_y \sin \theta). \end{aligned}$$

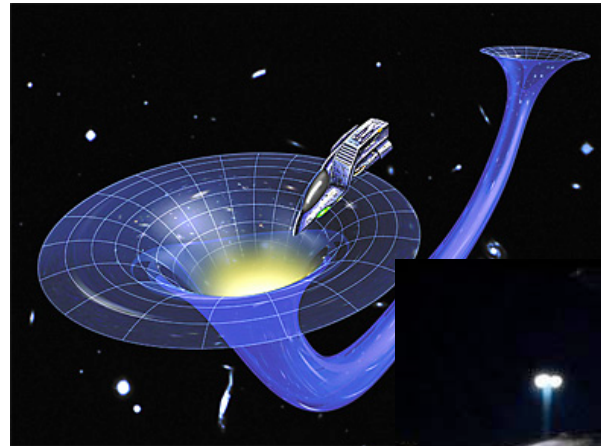
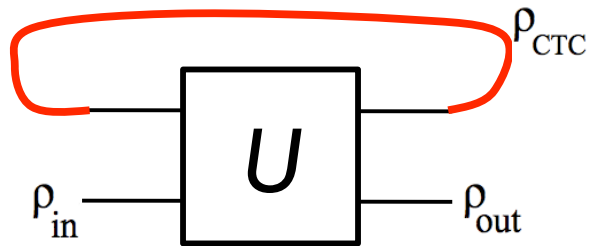
- Solution for generic input:

$$\begin{cases} \rho_{CTC} : \vec{m} = (n_z, 0, 0) \\ \rho_{out} : \vec{r} = (n_z^2, 0, 0) \end{cases} \leftarrow \boxed{\text{Non-linear map!}}$$

- Alternative solution for  $|\psi_{in}\rangle = \frac{1}{\sqrt{2}}(|0\rangle + e^{i\theta}|1\rangle)$  :  $\rho_{CTC} = \rho_{out} : \vec{m} = (0, 0, m_z)$



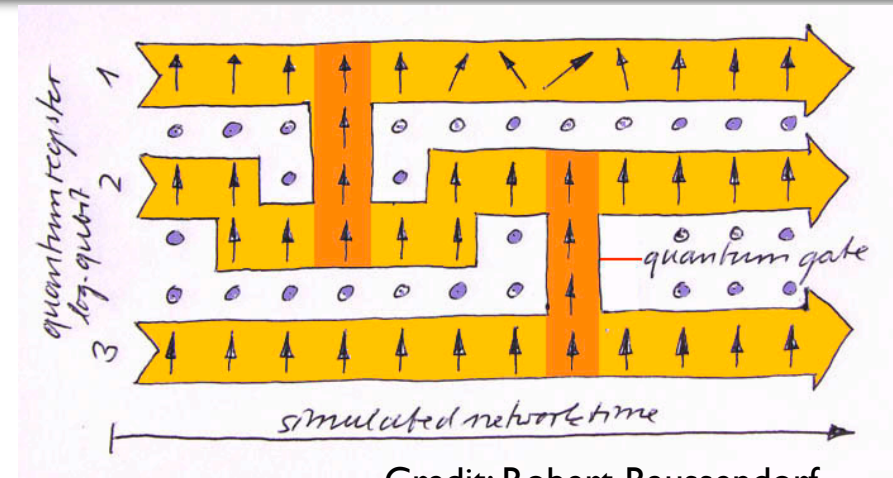
# Testing Deutsch's model



- Discussion of computational power of CTCs uses Deutsch's model. In the absence of experiments, how to check if the model is sound?
- We'll see that **measurement-based quantum computation** offers an answer.

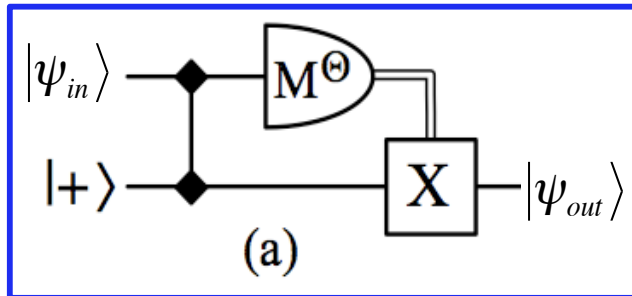
# The one-way model of quantum computing

- Proposed by Raussendorf/Briegel  
[PRL 86, 5188 (2001)]



- Consists in:
  - Preparation of standard entangled states via Heisenberg interactions – *cluster states*;
  - Adaptive sequence of 1-qubit measurements.
- Computational resource = quantum correlations of initial state
- Algorithm = choice of adaptive sequence of measurements

# Example: $J$ gate



$$|\psi_{in}\rangle = \alpha|0\rangle + \beta|1\rangle$$

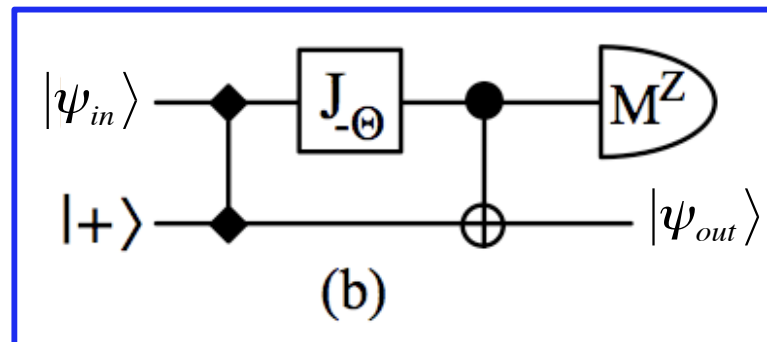
$$|+\rangle \equiv 1/\sqrt{2}(|0\rangle + |1\rangle)$$

$$M^\ominus = \text{Meas. on basis } \left\{ \begin{array}{l} |0\rangle + e^{i\theta}|1\rangle \leftrightarrow \text{outcome } s1 = 0 \\ |0\rangle - e^{i\theta}|1\rangle \leftrightarrow \text{outcome } s1 = 1 \end{array} \right\}$$

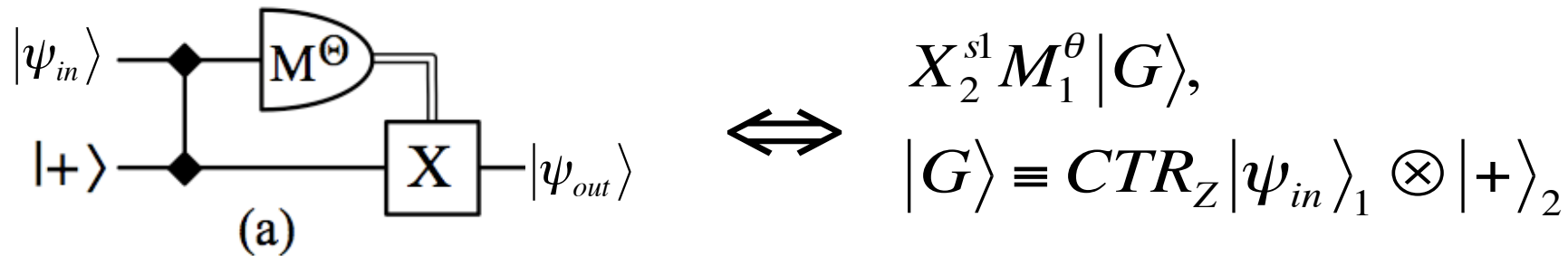
- Simple calculation shows that  $|\psi_{out}\rangle = J_{-\theta}|\psi_{in}\rangle$
- Circuit can be represented as command sequence:

$$X_2^{s1} M_1^\theta |G\rangle, \quad |G\rangle \equiv CTR_Z |\psi_{in}\rangle_1 \otimes |+\rangle_2$$

- Equivalent circuit – measure Z, implement  $CTR-X$  ( $CNOT$ ) coherently:



# J gate in MBQC = CTC



- Stabilizers of state  $|G\rangle$  :
 
$$\begin{cases} (Z_1 \otimes X_2)^0 = Z_1^0 \otimes X_2^0 = 1 \text{ (identity)} \\ (Z_1 \otimes X_2)^1 = Z_1^1 \otimes X_2^1 \end{cases}$$

That is,  $Z_1^{s_1} \otimes X_2^{s_1}$  is stabilizer independently of the outcome  $s_1$  of the measurement on qubit 1:

$$Z_1^{s_1} \otimes X_2^{s_1} |G\rangle = |G\rangle$$

- We can then perform stabilizer manipulation:

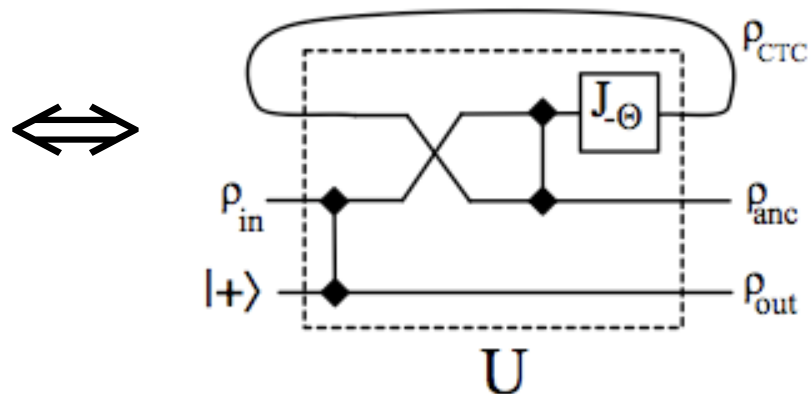
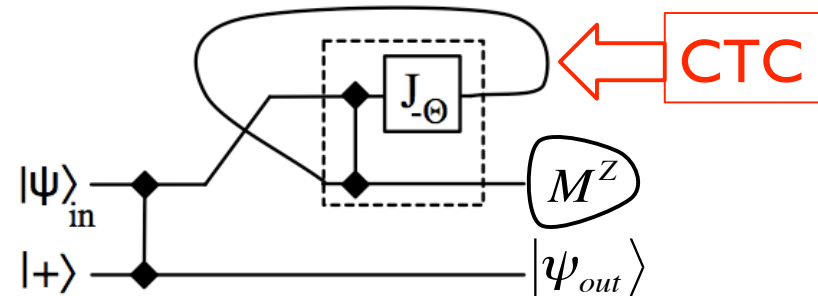
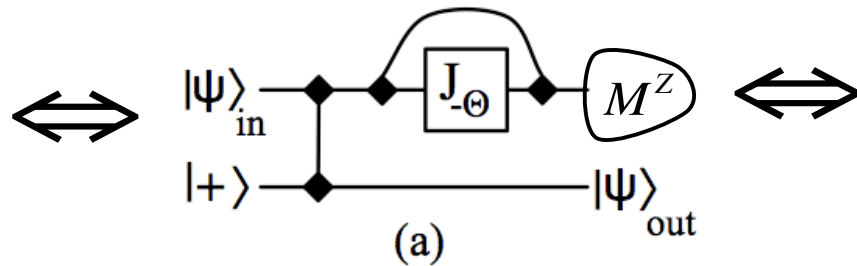
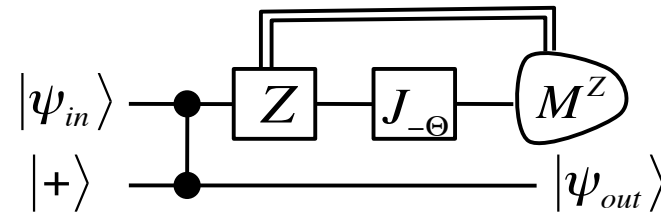
$$\begin{aligned} X_2^{s_1} M_1^\theta |G\rangle &= X_2^{s_1} M_1^\theta (Z_1^{s_1} X_2^{s_1} |G\rangle) \\ &= X_2^{s_1+s_1} M_1^\theta Z_1^{s_1} |G\rangle = \boxed{M_1^\theta Z_1^{s_1} |G\rangle} \end{aligned}$$

**Time-travel** situation: we need to apply Z depending on outcome of measurement not yet made.

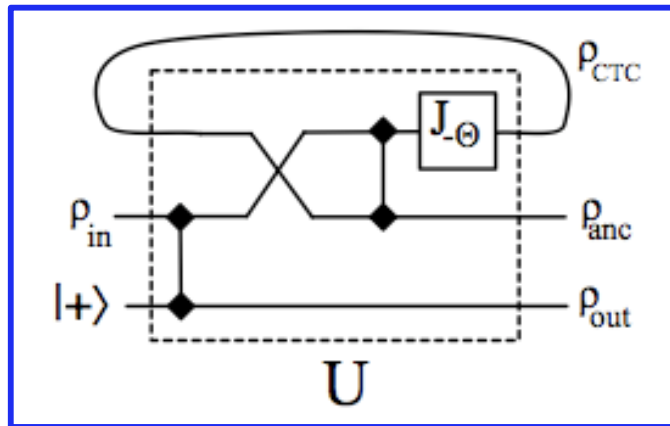
# $J$ gate in MBQC = CTC

- Putting this CTC in Deutsch format:

$$M_1^\theta Z_1^{s1} |G\rangle$$



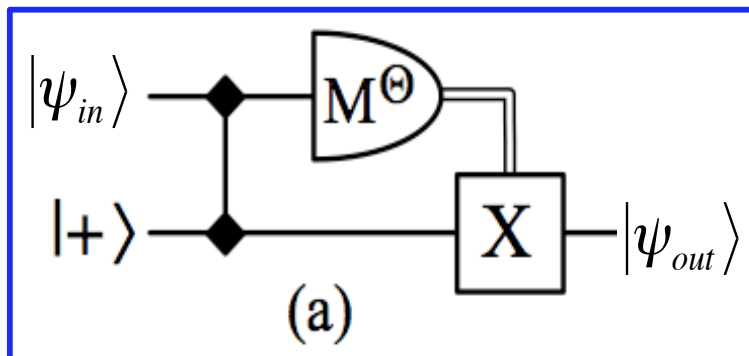
# Comparing the one-way model with Deutsch:



Deutsch:

$$\begin{cases} \rho_{in} : \vec{n} = (n_x, n_y, n_z) \\ \rho_{anc} : \vec{m} = (n_z^2, 0, 0) \\ \rho_{out} = \rho_{CTC} : \vec{r} = (n_z, 0, 0) \end{cases}$$

$$\rho_{in} = \frac{1}{2} \begin{pmatrix} 1 + n_z & n_x - in_y \\ n_x + in_y & 1 - n_z \end{pmatrix} \longrightarrow \rho_{out} = \frac{1}{2} \begin{pmatrix} 1 & n_z \\ n_z & 1 \end{pmatrix}$$



One-way model:

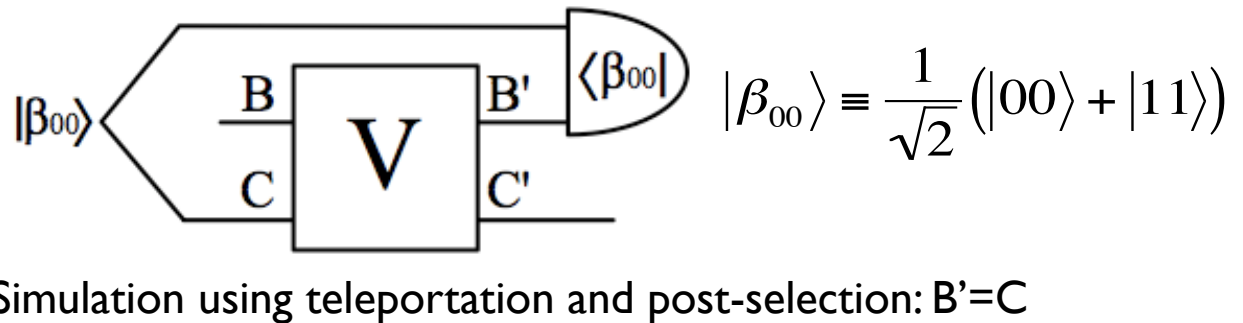
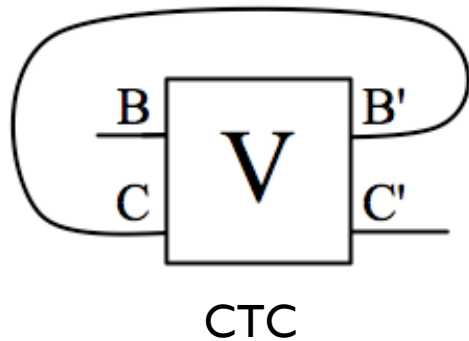
$$|\psi_{in}\rangle = \alpha|0\rangle + \beta|1\rangle$$

$$|\psi_{out}\rangle = J_{-\theta}|\psi_{in}\rangle = \alpha|+\rangle + \beta e^{-i\theta}|-\rangle$$

Deutsch's prediction is *different* from the one based on the one-way model. What's wrong with Deutsch?

# CTCs: model by Bennett/Schumacher/Svetlichny (BSS)

- Bennett and Schumacher, unpublished (2002) – see seminar <http://bit.ly/crs8Lb>
- Rediscovered independently by Svetlichny (2009) - [arXiv:0902.4898v1](https://arxiv.org/abs/0902.4898v1)
  - Related work on black holes by Horowitz/Maldacena (2004), Preskill/Gottesman (2004)

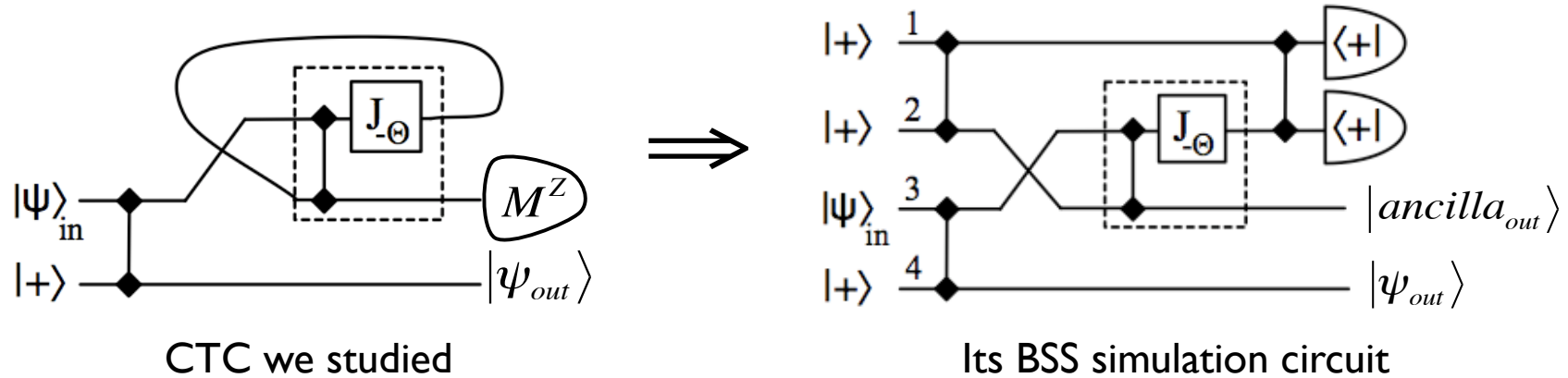


$$|\beta_{00}\rangle \equiv \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$$

- We post-select projections onto  $|\beta_{00}\rangle$ 
  - Postselection successful: state  $B'$  is teleported back in time (state  $C =$  state  $B'$ )
  - Simulation works only when post-selection happens -> finite probability of success.

What are BSS's predictions for our CTC?

# BSS x MBQC



- Simple calculation shows that BSS circuit implements (probabilistically) the map

$$|\psi_{in}\rangle \rightarrow |\psi_{out}\rangle = J_{-\theta} |\psi_{in}\rangle$$

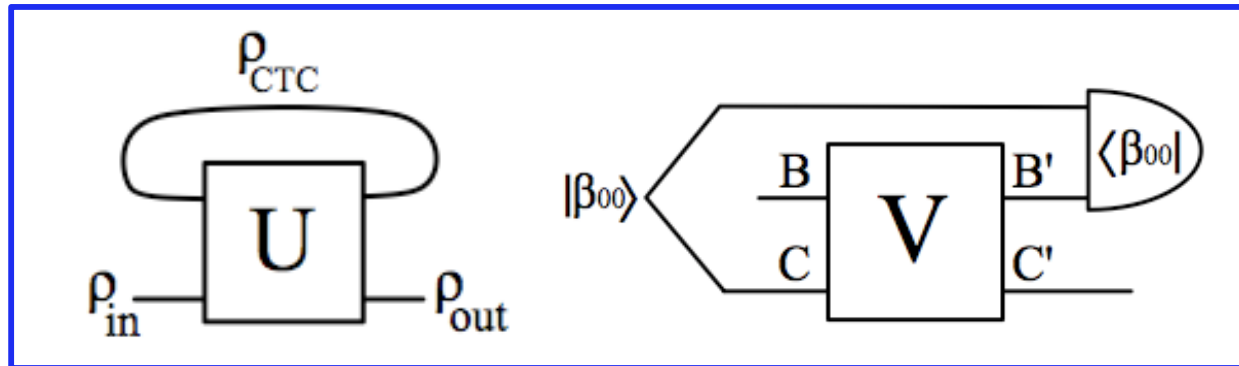
... recovering exactly MBQC's prediction!

**➔ BSS is the right model to explain CTCs in MBQC, and not Deutsch's...**



# Deutsch/BSS comparison

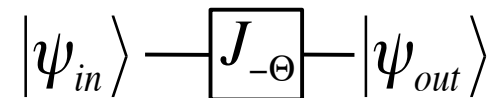
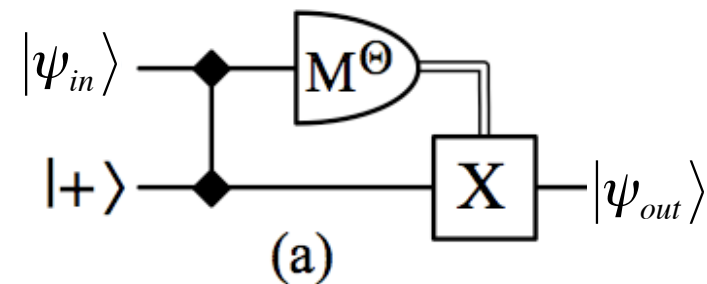
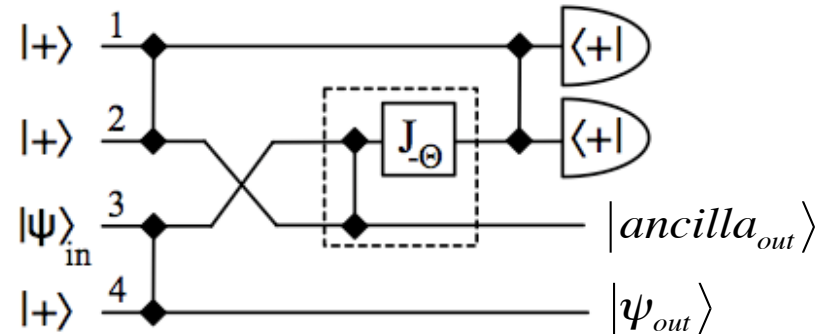
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- Deutsch's self-consistency is achieved with artificial mixed-states, CTC qubit sent back in time decorrelated with other systems.
- BSS preserves entanglement and correlations of CTC qubit due to teleportation step.
- Deutsch solutions with pure-state CTC qubit coincide with BSS solution.

# MBQC as deterministic simulations of CTCs

- Stabilizer techniques enable us to simplify BSS circuits
  - Z deletion;
  - Local complementation.
- Some BSS circuits reduce to MBQC patterns that deterministically simulate a unitary.
  - flow, gflow theorems.
- For example, the BSS circuit above is equivalent to:
- Or more simply:



# Conclusions

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- CTCs appear in MBQC; we analyze them using two CTC models.
- BSS's model **agrees** with MBQC.
- Deutsch's model is **in conflict** with MBQC – the CTC qubit is sent to past stripped of its entanglement.
- We characterize a class of CTCs that admit deterministic simulation circuits using the BSS model.
- More work is needed to better understand implications of the BSS model:
  - MQ + Deutsch's CTCs = PSPACE (Aaronson/Watrous 2008)
  - BSS is associated with complexity class PostBQP=PP (Aaronson 2004)
  - See recent work (and experiment) by Lloyd *et al.*: [arXiv:1005.2219v1](https://arxiv.org/abs/1005.2219v1)

# PP versus PSPACE

From [http://qwiki.stanford.edu/wiki/Complexity\\_Zoo](http://qwiki.stanford.edu/wiki/Complexity_Zoo)

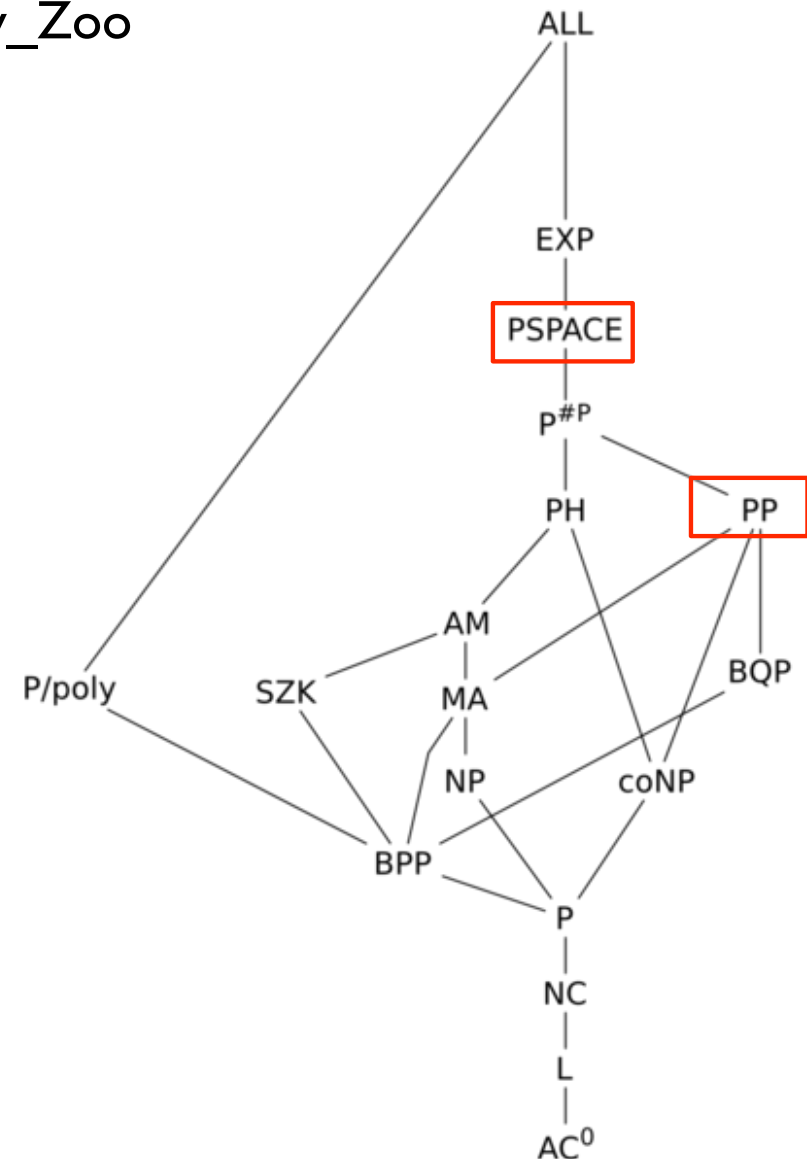
## PP

Like [BPP](#), [PP](#) is a class defined in an attempt to find out what randomness allows us to do algorithmically.

Formally, PP is the class of problems solvable by an [NP](#) machine such that, given a "yes" instance, strictly more than 1/2 of the computation paths accept, while given a "no" instance, strictly less than 1/2 of the computation paths accept

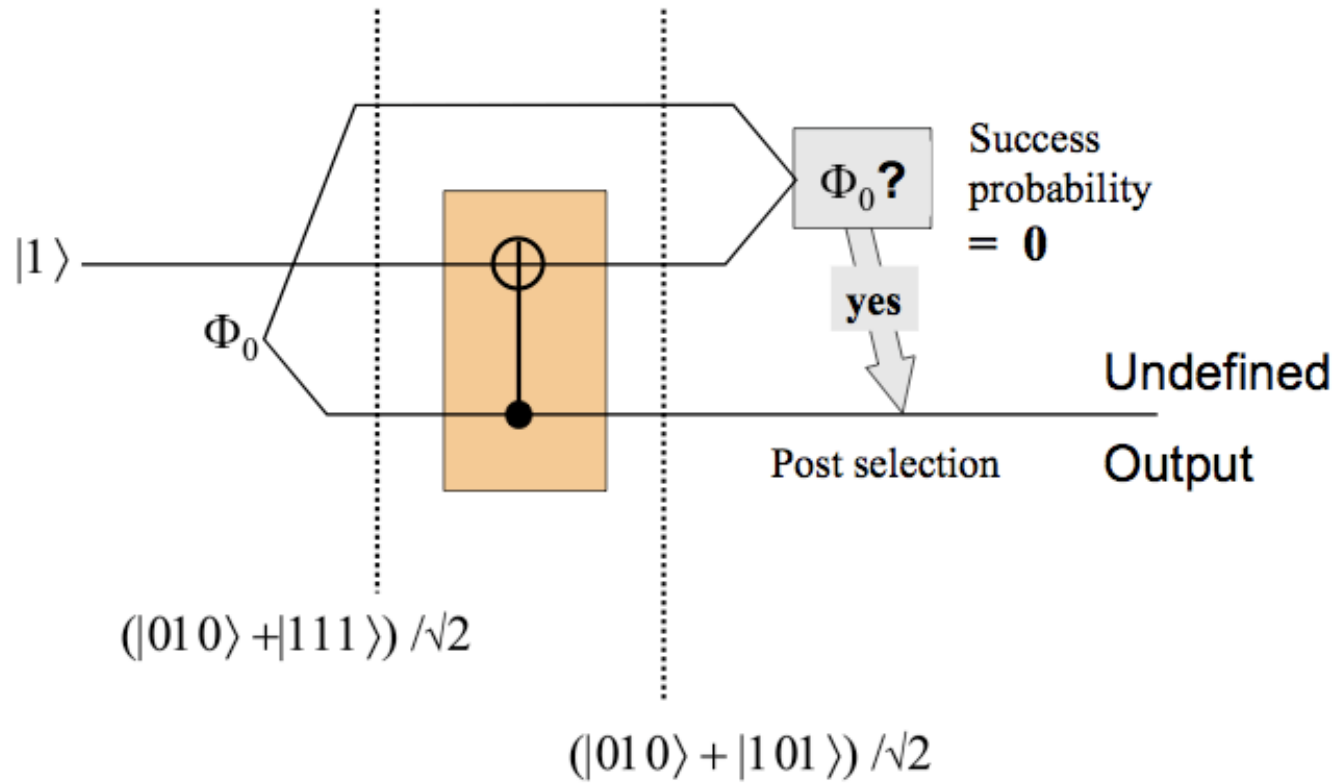
## PSPACE

Whereas [P](#) is a class of problems that can be solved in a polynomially-bounded amount of time, [PSPACE](#) is the class of problems that can be solved by a deterministic Turing machine that uses only a polynomially-bounded amount of space, regardless of how long the computation takes.



# How BSS deals with the grandfather paradox

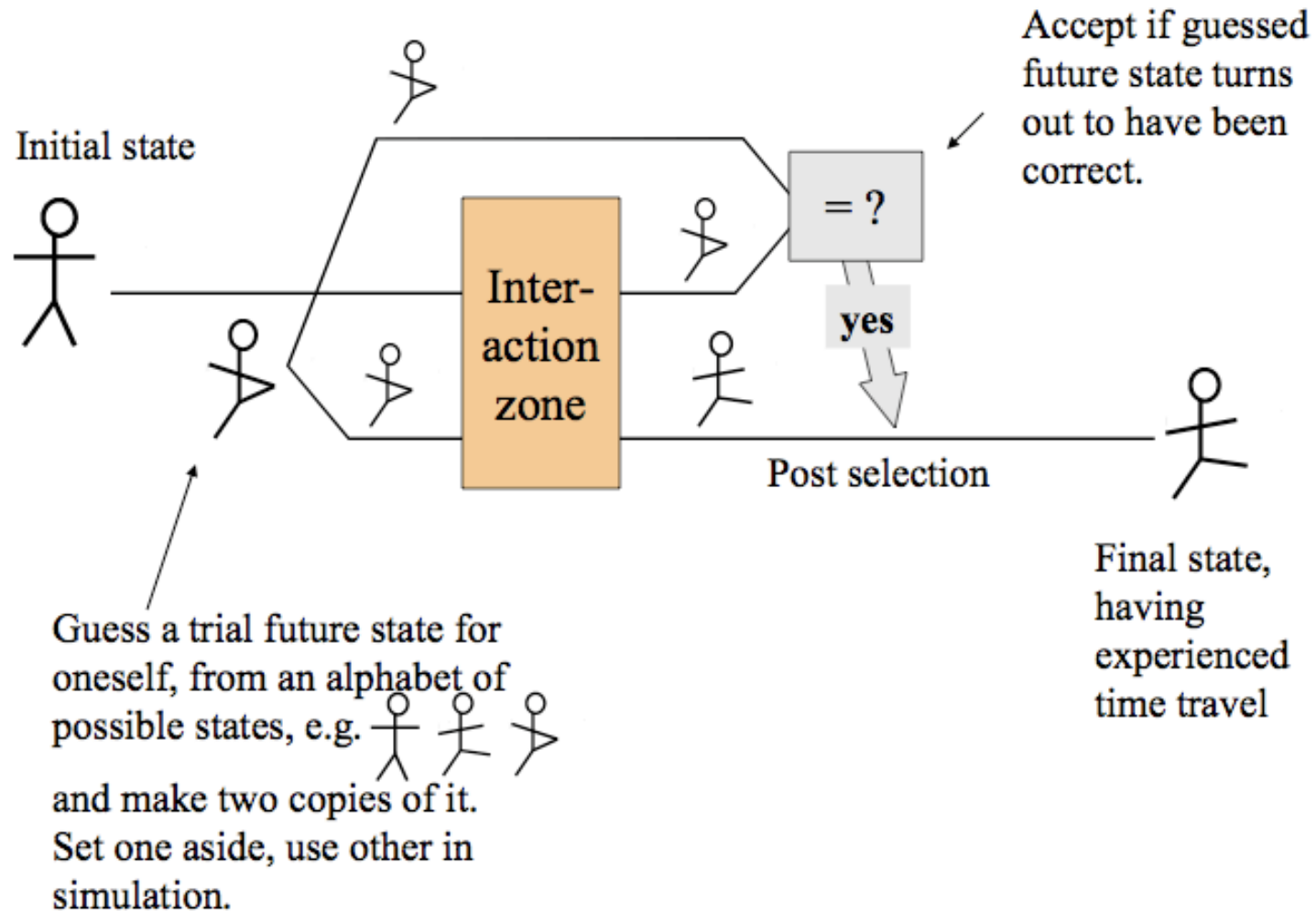
- From Bennett's talk slides: <http://bit.ly/crs8Lb>



- Paradoxical combinations of input and **unitary** result in post-selection with success probability  $p=0$ .

# Bennett's classical motivation for the BSS model

Woody Allen MC can be used to simulate time travel without need of any exotic physical equipment.



From Bennett's talk (2005).